FRED $9 / 2 / 71$ Binary Mode 3.


FRED $9 / 2 / 71$ Binay Mode $3^{\prime \prime}$
"FRED Lis a symbol stream processor. It takes as its input a stream of characters and produces as its output another stream of characters which is produced from the infut by direct copying except in the case of macro calls in the input stream, which are. evaluated before they are put into the output stream.

A macro call consists of a macro name and a list of parameters, each separated by a comma. The name is preceded by $*$ and the last pararneter followed by a semicolon.

$$
\text { e.g. } * M A C, 2,5 \text {; }
$$

Before the macro call can be evaluated the macro must have been defined by associating its name with a symbol string. This string may contain special symbols $\sim 1, \sim 2$, which stand for the first; second, etc, formal paraneters; the symbol $\sim 0$ stands for the cask name of the macro being evaluated.
e.g. if name $A B C$ defines string $A B \sim 1 C \sim 2 A B$ the call $* A B C, X Y, P Q$; will produce $A B X Y C P Q A B$

The system is completely general and it is possible to use a macro call in place of or in conjunction with a symbol string any-. where. In particular, macro calls are allowed in the actual parameters of other macro calls (including the name) and also in the defining string. The following examples demonstrate this point

| Name | Associated |
| :---: | :---: |
| APA | P~1~1P |
| $\begin{aligned} & \text { Macro-call } \\ & * A, C ; \\ & * A, * A, C ; \end{aligned}$ | Result <br> ACA <br> AACAA |
| ** $A, P ;, Y$; | PYyp |

Enclosing any string in the stringx quotes <...... $>$
has the effect of preventing evaluation of an macro calls inside; in place of an evaluation, however, one layer of string quotes is removed.

$$
\begin{aligned}
& \begin{array}{ll}
\text { e.g. Infut string } & \text { Result. } \\
Q\langle * A, C ;\rangle R & Q * A, C ; R
\end{array} \\
& Q\langle *\rangle R\langle;\rangle \\
& Q\langle\langle * A, C ;\rangle R \\
& \text { Q*R; } \\
& Q<* A, C ;>R
\end{aligned}
$$

The use of string quotes makes it possible to include any symbol in the output stream except an unmatched opening or closing string quote.

A macro is defined by a special macro DEF which has been writton in machino code mancluwed-in-the syatem. DEF takes two arguments: the name of the macro to be defined, and the defining symbol string. It is usual to enclose the symbol string in string quotes in order to prevent any mecro calls or uses of formal parameters from being effective during the process of definition. These quotes will be removed by the normal process of evaluating the arguments of DPF.

$$
\begin{aligned}
& * D E F, A,\langle A \sim \mid A\rangle ; \\
& * D E F, B,\langle B * A, X \sim \mid X ; B\rangle ; \\
& * D E F, A P A,\langle P \sim / \sim \mid P\rangle ;
\end{aligned}
$$

As definition is performed by an ordinary macro call the system insures that it is possible to carry out a definition anywhere it is possible to use a macro call. In particular a definition can be included in an actual parameter for a macro call and hence in the symbol string defining a macro.

In general the actual parameter list of a macro call is lost when the call hes been completed and this applies also to definitions that are part of the list. Definitions of this sort are therefore temporary and their scope is confined to this particuler macro call. If a macro name which has already been defined is defined again by . a call of DEF, the latest definition supersedes the earlier one, though without destroying it.

The basic macro UPDATE which takes two arguments has the same sort of effect as DEF except that instead of establishing a new definition it alters the value associated with its firstorargument to be its second argument. There is a limitation on the use of UPDATE as the space available for the value is fixed by the first definition, the new string may be of equal length or shorter.

Integer arithmetic is provided with the aid of three machine code macros: BIN converts a digit string, possibly preceded by a sign, into a signed binary integer.

DEC is the inverse operation, converting a signed binary interge into a decimal digit string of characters.

BAR takes three arguments, the first being the character + , - , ., /, or R , the other two being binary numbers. It performs the indiceted operation on these. BAR, $R, x, y$; gives the remainder when $x$ is divided by $y$.
Punching Rules:
$\circledast$
 not know about SiR subroutines.

CHIP roles for $920-903$ circe equivalents apply
Tapes may be punched in 920 code or 903 code and the first non-blank character on the tape must be a "Newline".

All tapes must and with an unmatched" $>$ " character followed by a "Newline". Stopeode is not recognised.
( $\star$
Sung ark:-

$$
\text { Cor 903, } 350 \text {, or ASCII coles winch in } 900 \text {-Serer telecurde }
$$ the character " to introduce format parameters,

B. lank, Erase, and Camaije seton ails

All lager must end cut t
? sorption of reit. reigned \& may same

Method of use.

1) Load FRED under initial instructions - Mode 3.
2) Put first tape in reader.
$\rightarrow$ (Frs output in 920 code tinges at 8)
swap. $\longrightarrow$ For output in aco-senies $\rightarrow 9-3$ code trigger at 10.
3) For subsequent tapes trigger at 14 .

Enos Indications.
If the input stream contanis contextual? errors an
*) error number will be output in legible tope, preceded "and brewed by $18^{\prime \prime}$ of blanks.
by -if Continuation after errors is not allowed. not andy not?
in error $\int$ Other tape errors such as party error, impimissille una? character, etc will cause a legible tape message of
(*) the form "CHI/O ERROR N" to be outport.
must list Contextual? ensor indications are:
1 - Unmatched ; in definition sting.
2 Unquoted $\sim$ in argument list.
3
4 Not enough arguments supplied in call
5 Probably a missing;-
6
7 Undefined macronavie.
8
9 Update string too long.

Also fees shoot than th manes is....

